

## 8257/8257-5 PROGRAMMABLE DMA CONTROLLER

- MCS-85<sup>®</sup> Compatible 8257-5
- 4-Channel DMA Controller
- Priority DMA Request Logic
- Channel Inhibit Logic
- Terminal Count and Modulo 128 Outputs

- Single TTL Clock
- Single + 5V Supply
- Auto Load Mode

## Available in EXPRESS Standard Temperature Range

The Intel<sup>®</sup> 8257 is a 4-channel direct memory access (DMA) controller. It is specifically designed to simplify the transfer of data at high speeds for the Intel<sup>®</sup> microcomputer systems. Its primary function is to generate, upon a peripheral request, a sequential memory address which will allow the peripheral to read or write data directly to or from memory. Acquisition of the system bus in accomplished via the CPU's hold function. The 8257 has priority logic that resolves the peripherals requests and issues a composite hold request to the CPU. It maintains the DMA cycle count for each channel and outputs a control signal to notify the peripheral that the programmed number of DMA cycles is complete. Other output control signals simplify sectored data transfers. The 8257 represents a significant savings in component count for DMA-based microcomputer systems and greatly simplifies the transfer of data at high speed between peripherals and memories.



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Figure 1. Block Diagram

Figure 2. Pin Configuration

## FUNCTIONAL DESCRIPTION

#### General

The 8257 is a programmable, Direct Memory Access (DMA) device which, when coupled with a single Intel<sup>®</sup> 8212 I/O port device, provides a complete four-channel DMA controller for use in Intel<sup>®</sup> microcomputer systems. After being initialized by software, the 8257 can transfer a block of data, containing up to 16,384 bytes, between memory and a peripheral device directly, without further intervention required of the CPU. Upon receiving a DMA transfer request from an enabled peripheral, the 8257:

- 1. Acquires control of the system bus.
- 2. Acknowledges that requesting peripheral which is connected to the highest priority channel.
- 3. Outputs the least significant eight bits of the memory address onto system address lines A<sub>0</sub>-A<sub>7</sub>, outputs the most significant eight bits of the memory address to the 8212 I/O port via the data bus (the 8212 places these address bits on lines A<sub>8</sub>-A<sub>15</sub>), and
- 4. Generates the appropriate memory and I/O read/ write control signals that cause the peripheral to receive or deposit a data byte directly from or to the addressed location in memory.

The 8257 will retain control of the system bus and repeat the transfer sequence, as long as a peripheral maintains its DMA request. Thus, the 8257 can transfer a block of data to/from a high speed peripheral (e.g., a sector of data on a floppy disk) in a single "burst". When the specified number of data bytes have been transferred, the 8257 activates its Terminal Count (TC) output, informing the CPU that the operation is complete.

The 8257 offers three different modes of operation: (1) DMA read, which causes data to be transferred from memory to a peripheral; (2) DMA write, which causes data to be transferred from a peripheral to memory; and (3) DMA verify, which does not actually involve the transfer of data. When an 8257 channel is in the DMA verify mode, it will respond the same as described for transfer operations, except that no memory or I/O read/write control signals will be generated, thus preventing the transfer of data. The 8257, however, will gain control of the system bus and will acknowledge the peripheral's DMA request for each DMA cycle. The peripheral can use these acknowledge signals to enable an internal access of each byte of a data block in order to execute some verification procedure, such as the accumulation of a CRC (Cyclic Redundancy Code) checkword. For example, a block of DMA verify cycles might follow a block of DMA read cycles (memory to peripheral) to allow the peripheral to verify its newly acquired data.

#### **Block Diagram Description**

#### 1. DMA Channels

The 8257 provides four separate DMA channels (labeled CH-0 to CH-3). Each channel includes two sixteen-bit registers: (1) a DMA address register, and (2) a terminal count register. Both registers must be initialized before a channel is enabled. The DMA address register is loaded with the address of the first memory location to be accessed. The value loaded into the low-order 14-bits of the terminal count register specifies the number of DMA cycles minus one before the Terminal Count (TC) output is activated. For instance, a terminal count of 0 would cause the TC output to be active in the first DMA cycle for that channel. In general, if N = the number of desired DMA cycles, load the value N-1 into the low-order 14-bits of the terminal count register. The most significant two bits of the terminal count register specify the type of DMA operation for that channel.



Figure 3. 8257 Block Diagram Showing DMA Channels

These two bits are not modified during a DMA cycle, but can be changed between DMA blocks.

Each channel accepts a DMA Request (DRQn) input and provides a DMA Acknowledge (DACKn) output.

#### (DRQ 0-DRQ 3)

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DMA Request: These are individual asynchronous channel request inputs used by the peripherals to obtain a DMA cycle. If not in the rotating priority mode then DRQ 0 has the highest priority and DRQ 3 has the lowest. A request can be generated by raising the request line and holding it high until DMA acknowledge. For multiple DMA cycles (Burst Mode) the request line is held high until the DMA acknowledge of the last cycle arrives.

#### (DACK 0 - DACK 3)

DMA Acknowledge: An active low level on the acknowledge output informs the peripheral connected to that channel that it has been selected for a DMA cycle. The DACK output acts as a "chip select" for the peripheral device requesting service. This line goes active (low) and inactive (high) once for each byte transferred even if a burst of data is being transferred.

#### 2. Data Bus Buffer

This three-state, bi-directional, eight bit buffer interfaces the 8257 to the system data bus.

#### $(D_0 - D_7)$

Data Bus Lines: These are bi-directional three-state lines. When the 8257 is being programmed by the CPU, eightbits of data for a DMA address register, a terminal count register or the Mode Set register are received on the data bus. When the CPU reads a DMA address register, a terminal count register or the Status register, the data is sent to the CPU over the data bus. During DMA cycles (when the 8257 is the bus master), the 8257 will output the most significant eight-bits of the memory address (from one of the DMA address registers) to the 8212 latch via the data bus. These address bits will be transferred at the beginning of the DMA cycle; the bus will then be released to handle the memory data transfer during the balance of the DMA cycle.

<b>BIT</b> 15	BIT 14	TYPE OF DMA OPERATION
0	0	Verify DMA Cycle
0	1	Write DMA Cycle
1	0	Read DMA Cycle
1	1	(Illegal)



Figure 4. 8257 Block Diagram Showing Data Bus Buffer

#### 3. Read/Write Logic

When the CPU is programming or reading one of the 8257's registers (i.e., when the 8257 is a "slave" device on the system bus), the Read/Write Logic accepts the I/O Read (I/OR) or I/O Write (I/OW) signal, decodes the least significant four address bits,  $(A_0-A_3)$ , and either writes the contents of the data bus into the addressed register (if I/OW is true) or places the contents of the addressed register onto the data bus (if I/OR is true).

During DMA cycles (i.e., when the 8257 is the bus "master"), the Read/Write Logic generates the I/O read and memory write (DMA write cycle) or I/O Write and memory read (DMA read cycle) signals which control the data link with the peripheral that has been granted the DMA cycle.

Note that during DMA transfers Non-DMA I/O devices should be de-selected (disabled) using "AEN" signal to inhibit I/O device decoding of the memory address as an erroneous device address.

#### (I/OR)

I/O Read: An active-low, bi-directional three-state line. In the "slave" mode, it is an input which allows the 8-bit status register or the upper/lower byte of a 16-bit DMA address register or terminal count register to be read. In the "master" mode, I/OR is a control output which is used to access data from a peripheral during the DMA write cycle.

#### (1/OW)

I/O Write: An active-low, bi-directional three-state line. In the "slave" mode, it is an input which allows the contents of the data bus to be loaded into the 8-bit mode set register or the upper/lower byte of a 16-bit DMA address register or terminal count register. In the "master" mode, I/OW is a control output which allows data to be output to a peripheral during a DMA read cycle.

#### (CLK)

Clock Input: Generally from an Intel® 8224 Clock Generator device. ( $\phi$ 2 TTL) or Intel® 8085A CLK output.

#### (RESET)

Reset: An asynchronous input (generally from an 8224 or 8085 device) which disables all DMA channels by clearing the mode register and 3-states all control lines.

#### (A<sub>0</sub>-A<sub>3</sub>)

Address Lines: These least significant four address lines are bi-directional. In the "slave" mode they are inputs which select one of the registers to be read or programmed. In the "master" mode, they are outputs which constitute the least significant four bits of the 16-bit memory address generated by the 8257.

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#### (CS)

Chip Select: An active-low input which enables the I/O Read or I/O Write input when the 8257 is being read or programmed in the "slave" mode. In the "master" mode, CS is automatically disabled to prevent the chip from selecting itself while performing the DMA function.

#### 4. Control Logic

This block controls the sequence of operations during all DMA cycles by generating the appropriate control signals and the 16-bit address that specifies the memory location to be accessed.



Figure 5. 8257 Block Diagram Showing -Read/Write Logic Function

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#### (A4-A7)

Address Lines: These four address lines are three-state outputs which constitute bits 4 through 7 of the 16-bit memory address generated by the 8257 during all DMA cycles.

#### (READY)

Ready: This asynchronous input is used to elongate the memory read and write cycles in the 8257 with wait states if the selected memory requires longer cycles. READY must conform to specified setup and hold times.

#### (HRQ)

Hold Request: This output requests control of the system bus. In systems with only one 8257, HRQ will normally be applied to the HOLD input on the CPU. HRQ must conform to specified setup and hold times.

#### (HLDA)

Hold Acknowledge: This input from the CPU indicates that the 8257 has acquired control of the system bus.

#### (MEMR)

Memory Read: This active-low three-state output is used to read data from the addressed memory location during DMA Read cycles.

#### (MEMW)

Memory Write: This active-low three-state output is used to write data into the addressed memory location during DMA Write cycles.

#### (ADSTB)

Address Strobe: This output strobes the most significant byte of the memory address into the 8212 device from the data bus.

#### (AEN)

Address Enable: This output is used to disable (float) the System Data Bus and the System Control Bus. It may also be used to disable (float) the System Address Bus by use of an enable on the Address Bus drivers in systems to inhibit non-DMA devices from responding during DMA cycles. It may be further used to isolate the 8257 data bus from the System Data Bus to facilitate the transfer of the 8 most significant DMA address bits over the 8257 data I/O pins without subjecting the System Data Bus to any timing constraints for the transfer. When the 8257 is used in an I/O device structure (as opposed to memory mapped), this AEN output should be used to disable the selection of an I/O device when the DMA address is on the address bus. The I/O device selection should be determined by the DMA acknowledge outputs for the 4 channels.

#### (TC)

Terminal Count: This output notifies the currently selected peripheral that the present DMA cycle should be the last cycle for this data block. If the TC STOP bit in the Mode Set register is set, the selected channel will be automatically disabled at the end of that DMA cycle. TC is activated when the 14-bit value in the selected channel's terminal count register equals zero. Recall that the low-order 14-bits of the terminal count register should be loaded with the values (n-1), where n = the desired number of the DMA cycles.

#### (MARK)

Modulo 128 Mark: This output notifies the selected peripheral that the current DMA cycle is the 128th cycle since the previous MARK output. MARK always occurs at 128 (and all multiples of 128) cycles from the end of the data block. Only if the total number of DMA cycles (n) is evenly divisable by 128 (and the terminal count register was loaded with n-1), will MARK occur at 128 (and each succeeding multiple of 128) cycles from the beginning of the data block.



Figure 6. 8257 Block Diagram Showing Control Logic and Mode Set Register

#### 5. Mode Set Register

When set, the various bits in the Mode Set register enable each of the four DMA channels, and allow four different options for the 8257:



The Mode Set register is normally programmed by the CPU after the DMA address register(s) and terminal count register(s) are initialized. The Mode Set Register is cleared by the RESET input, thus disabling all options, inhibiting all channels, and preventing bus conflicts on power-up. A channel should not be left enabled unless its DMA address and terminal count registers contain valid values; otherwise, an inadvertent DMA request (DRQn) from a peripheral could initiate a DMA cycle that would destroy memory data.

The various options which can be enabled by bits in the Mode Set register are explained below:

#### **Rotating Priority Bit 4**

In the Rotating Priority Mode, the priority of the channels has a circular sequence. After each DMA cycle, the priority of each channel changes. The channel which had just been serviced will have the lowest priority.



If the ROTATING PRIORITY bit is not set (set to a zero), each DMA channel has a fixed priority. In the fixed priority mode, Channel 0 has the highest priority and Channel 3 has the lowest priority. If the ROTATING PRIORITY bit is set to a one, the priority of each channel changes after each DMA cycle (not each DMA request). Each channel moves up to the next highest priority assignment, while the channel which has just been serviced moves to the lowest priority assignment:

	CHANNEL->	СН-0	СН-1	СН-2	СН-3
Priority	Highest	CH-1	CH-2	СН-3	CH-0
Assignments	<b>A</b>	CH-2	CH-3	CH-0	CH-1
	★	CH-3	CH-0	CH-1	CH-2
	Lowest	CH-0	CH-1	CH-2	СН-3

Note that rotating priority will prevent any one channel from monopolizing the DMA mode; consecutive DMA cycles will service different channels if more than one channel is enabled and requesting service. There is no overhead penalty associated with this mode of operation. All DMA operations began with Channel 0 initially assigned to the highest priority for the first DMA cycle.

#### **Extended Write Bit 5**

If the EXTENDED WRITE bit is set, the duration of both the MEMW and I/OW signals is extended by activating them earlier in the DMA cycle. Data transfers within microcomputer systems proceed asynchronously to allow use of various types of memory and I/O devices with different access times. If a device cannot be accessed within a specific amount of time it returns a "not ready" indication to the 8257 that causes the 8257 to insert one or more wait states in its internal sequencing. Some devices are fast enough to be accessed without the use of wait states, but if they generate their READY response with the leading edge of the I/OW or MEMW signal (which generally occurs late in the transfer sequence), they would normally cause the 8257 to enter a wait state because it does not receive READY in time. For systems with these types of devices, the Extended Write option provides alternative timing for the I/O and memory write signals which allows the devices to return an early READY and prevents the unnecessary occurrence of wait states in the 8257, thus increasing system throughput.

#### TC Stop Bit 6

If the TC STOP bit is set, a channel is disabled (i.e., its enable bit is reset) after the Terminal Count (TC) output goes true, thus automatically preventing further DMA operation on that channel. The enable bit for that channel must be re-programmed to continue or begin another DMA operation. If the TC STOP bit is not set, the occurrence of the TC output has no effect on the channel enable bits. In this case, it is generally the responsibility of the peripheral to cease DMA requests in order to terminate a DMA operation.

#### Auto Load Bit 7

The Auto Load mode permits Channel 2 to be used for repeat block or block chaining operations, without immediate software intervention between blocks. Channel 2 registers are initialized as usual for the first data block: Channel 3 registers, however, are used to store the block re-initialization parameters (DMA starting address, terminal count and DMA transfer mode). After the first block of DMA cycles is executed by Channel 2 (i.e., after the TC output goes true), the parameters stored in the Channel 3 registers are transferred to Channel 2 during an "update" cycle. Note that the TC STOP feature, described above, has no effect on Channel 2 when the Auto Load bit is set. If the Auto Load bit is set, the initial parameters for Channel 2 are automatically duplicated in the Channel 3 registers when Channel 2 is programmed. This permits repeat block operations to be set up with the programming of a single channel. Repeat block operations can be used in applications such as CRT refreshing. Channels 2 and 3 can still be loaded with separate values if Channel 2 is loaded before loading Channel 3. Note that in the Auto Load mode, Channel 3 is still available to the user if the Channel 3 enable bit is set, but use of this channel will change the values to be auto loaded into Channel 2 at update time. All that is necessary to use the Auto Load feature for chaining operations is to reload Channel 3 registers at the conclusion of each update cycle with the new parameters for the next data block transfer.

Each time that the 8257 enters an update cycle, the update flag in the status register is set and parameters in Channel 3 are transferred to Channel 2, non-destructively for Channel 3. The actual re-initialization of Channel 2 occurs at the beginning of the next channel 2 DMA cycle after the TC cycle. This will be the first DMA cycle of the new data block for Channel 2. The update flag is cleared at the conclusion of this DMA cycle. For chaining operations, the update flag in the status register can be monitored by the CPU to determine when the re-initialization process has been completed so that the next block parameters can be safely loaded into Channel 3.

#### 6. Status Register

The eight-bit status register indicates which channels have reached a terminal count condition and includes the update flag described previously.



The TC status bits are set when the Terminal Count (TC) output is activated for that channel. These bits remain set until the status register is read or the 8257 is reset. The UPDATE FLAG, however, is not affected by a status register read operation. The UPDATE FLAG can be cleared by resetting the 8257, by changing to the non-auto load mode (i.e., by resetting the AUTO LOAD bit in the Mode Set register) or it can be left to clear itself at the completion of the update cycle. The purpose of the UPDATE FLAG is to prevent the CPU from inadvertently skipping a data block by overwriting a starting address or terminal count in the Channel 3 registers before those parameters are properly auto-loaded into Channel 2.

The user is cautioned against reading the TC status register and using this information to reenable channels that have not completed operation. Unless the DMA channels are inhibited a channel could reach terminal count (TC) between the status read and the mode write. DMA can be inhibited by a hardware gate on the HRQ line or by disabling channels with a mode word before reading the TC status.



Figure 7. Autoload Timing

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## **OPERATIONAL SUMMARY**

#### Programming and Reading the 8257 Registers

There are four pairs of "channel registers": each pair consisting of a 16-bit DMA address register and a 16-bit terminal count register (one pair for each channel). The 8257 also includes two "general registers": one 8-bit Mode Set register and one 8-bit Status register. The registers are loaded or read when the CPU executes a write or read instruction that addresses the 8257 device and the appropriate register within the 8257. The 8228 generates the appropriate read or write control signal (generally I/OR or I/OW while the CPU places a 16-bit address on the system address bus, and either outputs the data to be written onto the system data bus or accepts the data being read from the data bus. All or some of the most significant 12 address bits A4-A15 (depending on the systems memory, I/O configuration) are usually decoded to produce the chip select ( $\overline{CS}$ ) input to the 8257. An I/O Write input (or Memory Write in memory mapped I/O configurations, described below) specifies that the addressed register is to be programmed, while an I/O Read input (or Memory Read) specifies that the addressed register is to be read. Address bit 3 specifies whether a "channel register" ( $A_3 = 0$ ) or the Mode Set (program only)/Status (read only) register  $(A_3 = 1)$  is to be accessed.

The least significant three address bits,  $A_0$ - $A_2$ , indicate the specific register to be accessed. When accessing the Mode Set or Status register,  $A_0$ - $A_2$  are all zero. When accessing a channel register bit  $A_0$  differentiates between the DMA address register ( $A_0 = 0$ ) and the terminal count register ( $A_0 = 1$ ), while bits  $A_1$  and  $A_2$  specify one of the

CONTROL INPUT	cs	I/OW	I/OR	CA
Program Haif of a Channel Register	0	0	1	0
Read Half of a Channel Register	0	1	0	0
Program Mode Set Register	0	0	1	1
<b>Read Status Register</b>	0	1	0	1

four channels. Because the "channel registers" are 16bits, two program instruction cycles are required to load or read an entire register. The 8257 contains a first/last (F/L) flip flop which toggles at the completion of each channel program or read operation. The F/L flip flop determines whether the upper or lower byte of the register is to be accessed. The F/L flip flop is reset by the RESET input and whenever the Mode Set register is loaded. To maintain proper synchronization when accessing the "channel registers" all channel command instruction operations should occur in pairs, with the lower byte of a register always being accessed first. Do not allow  $\overline{CS}$  to clock while either 1/OR or 1/OW is active, as this will cause an erroneous F/L flip flop state. In systems utilizing an interrupt structure, interrupts should be disabled prior to any paired programming operations to prevent an interrupt from splitting them. The result of such a split would leave the F/L F/F in the wrong state. This problem is particularly obvious when other DMA channels are programmed by an interrupt structure.

	1	AD	DRES	S INPL	JTS			.BI-	DIREC	TION	AL DA	TA BL	JS	
REGISTER	BYTE	A <sub>3</sub>	<b>A</b> 2	<b>A</b> 1	A <sub>0</sub>	F/L	D7	D <sub>6</sub>	D5	D4	D3	D <sub>2</sub>	D <sub>1</sub>	Do
CH-0 DMA Address	LSB MSB	0	0	0	0	0	A7 A15	A <sub>6</sub> A <sub>14</sub>	A5 A13	A4 A12	A <sub>3</sub> A <sub>11</sub>	A <sub>2</sub> A <sub>10</sub>	A <sub>1</sub> A <sub>9</sub>	A <sub>0</sub> A6
CH-0 Terminal Count	LSB MSB	0 0	0	0	1	0 1	C <sub>7</sub> Rd	C <sub>6</sub> Wr	C5 C13	C4 C12	C3 C11	C <sub>2</sub> C <sub>10</sub>	C <sub>1</sub> C <sub>9</sub>	C <sub>0</sub> C <sub>8</sub>
CH-1 DMA Address	LSB MSB	0	0	1	0	0 1	Same	as Cha	innel (					
CH-1 Terminal Count	LSB MSB	0 0	0	1	1	0 1								
CH-2 DMA Address	LSB MSB	0 0	1	0	0	0	Same	as Cha I	i nnel (	)				
CH-2 Terminal Count	LSB MSB	0	1	0	1	0								
CH-3 DMA Address	LSB MSB	0	1	1	0	0	Same	 as Chi 	 nnel ( 					
CH-3 Terminal Count	LSB MSB	0.0	1	1	1	0								
MODE SET (Program only)	-	1	0	0	0	0	AL	TCS	EW	RP	EN3	EN2	EN1	ENO
STATUS (Read only)	-	1	0	0	0	0	0	0	0	UP	тсз	TC2	TC1	TCO

#### 8257 Register Selection

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\*A0-A15: DMA Starting Address, C0-C13: Terminal Count value (N-1), Rd and Wr: DMA Verify (00), Write (01) or Read (10) cycle selection, AL: Auto Load, TCS: TC STOP, EW: EXTENDED WRITE, RP: ROTATING PRIORITY, EN3-EN0: CHANNEL ENABLE MASK, UP: UPDATE FLAG, TC3-TC0: TERMINAL COUNT STATUS BITS. 8257/8257-5

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Figure 8. DMA Operation State Diagram

#### DMA OPERATION

#### Singlé Byte Transfers

A single byte transfer is initiated by the I/O device raising the DRQ line of one channel of the 8257. If the channel is enabled, the 8257 will output a HRQ to the CPU. The 8257 now waits until a HLDA is received insuring that the system bus is free for its use. Once HLDA is received the DACK line for the requesting channel is activated (LOW). The DACK line acts as a chip select for the requesting I/O device. The 8257 then generates the read and write commands and byte transfer occurs between the selected I/O device and memory. After the transfer is complete, the DACK line is set HIGH and the HRQ line is set LOW to indicate to the CPU that the bus is now free for use. DRQ must remain HIGH until DACK is issued to be recognized and must go LOW before S4 of the transfer sequence to prevent another transfer from occuring. (See timing diagram.)

#### **Consecutive Transfers**

If more than one channel requests service simultaneously, the transfer will occur in the same way a burst does. No overhead is incurred by switching from one channel to another. In each S4 the DRQ lines are sampled and the highest priority request is recognized during the next transfer. A burst mode transfer in a lower priority channel will be overridden by a higher priority request. Once the high priority transfer has completed control will return to the lower priority channel if its DRQ is still active. No extra cycles are needed to execute this sequence and the HRQ line remains active until all DRQ lines go LOW.

#### **Control Override**

The continuous DMA transfer mode described above can be interrupted by an external device by lowering the HLDA line. After each DMA transfer the 8257 samples the HLDA line to insure that it is still active. If it is not active, the 8257 completes the current transfer, releases the HRQ line (LOW) and returns to the idle state. If DRQ lines are still active the 8257 will raise the HRQ line in the third cycle and proceed normally. (See timing diagram.)

#### **Not Ready**

The 8257 has a Ready input similar to the 8080A and the 8085A. The Ready line is sampled in State 3. If Ready is LOW the 8257 enters a wait state. Ready is sampled during every wait state. When Ready returns HIGH the 8257 proceeds to State 4 to complete the transfer. Ready is used to interface memory or I/O devices that cannot meet the bus set up times required by the 8257.

#### Speed

The 8257 uses four clock cycles to transfer a byte of data. No cycles are lost in the master to master transfer maximizing bus efficiency. A 2MHz clock input will allow the 8257 to transfer at a rate of 500K bytes/second.

#### **Memory Mapped I/O Configurations**

The 8257 can be connected to the system bus as a memory device instead of as an I/O device for memory mapped I/O configurations by connecting the system memory control lines to the 8257's 1/O control lines and the system I/O control lines to the 8257's memory control lines.

This configuration permits use of the 8080's considerably larger repertoire of memory instructions when reading or loading the 8257's registers. Note that with this connection, the programming of the Read (bit 15) and Write (bit 14) bits in the terminal count register will have a different meaning: intel

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BIT 15 READ	BIT 14 WRITE	
0	0,	DMA Verify Cycle
0	1	DMA Read Cycle
.1	0	DMA Write Cycle
1	• 1	lliegal

Figure 9. System Interface for Memory Mapped I/O

Figure 10. TC Register for Memory Mapped I/O Only

## SYSTEM APPLICATION EXAMPLES



Figure 11. Floppy Disk Controller (4 Drives)



Figure 12. High-Speed Communication Controller 10. 101

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A.C. TESTING LOAD CIRCUIT



## A.C. TESTING INPUT, OUTPUT WAVEFORM



#### **Tracking Parameters**

Signals labeled as Tracking Parameters (footnotes 1 and 5-7 under A.C. Specifications) are signals that follow similar paths through the silicon die. The propagation speed of these signals varies in the manufacturing process but the relationship between all these parameters is constant. The variation is less than or equal to 50 ns.

Suppose the following timing equation is being evaluated,

 $T_{A(MIN)} + T_{B(MAX)} \le 150 \text{ ns}$ 

and only minimum specifications exist for  $T_A$  and  $T_B$ . If  $T_{A(MIN)}$  is used, and if  $T_A$  and  $T_B$  are tracking parameters,  $T_{B(MAX)}$  can be taken as  $T_{B(MIN)}$  + 50 ns.

 $T_{A(MIN)} + (T_{B(MIN)}^{*} + 50 \text{ ns}) \le 150 \text{ ns}$ 

\*if  $T_A$  and  $T_B$  are tracking parameters

#### WAVEFORMS—PERIPHERAL MODE



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## **WAVEFORMS (Continued)**





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Figure 13. Detailed System Interface Schematic

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## ABSOLUTE MAXIMUM RATINGS\*

Ambient Temperature Under Bias .	0°C to 70°C
Storage Temperature	65°C to +150°C
Voltage on Any Pin	•
With Respect to Ground	0.5V to +7V
Power Dissipation	1 Watt

\*NOTICE: Stresses above those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

## **D.C. CHARACTERISTICS**

(8257:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 5\%$ , GND = 0V) (8257-5:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 10\%$ , GND = 0V)

Symbol	Parameter	Min.	Max.	Unit	Test Conditions
VIL	Input Low Voltage	-0.5	0.8	Volts	
ViH	Input High Voltage	2.0	V <sub>CC</sub> +.5	Volts	
VOL	Output Low Voltage		0.45	Volts	l <sub>OL</sub> = 1.6 mA
Voн	Output High Voltage	2.4	Vcc	Volts	$l_{OH} = -150 \mu A$ for AB, DB and AEN $l_{OH} = -80 \mu A$ for others
V <sub>HH</sub>	HRQ Output High Voltage	3.3	Vcc	Volts	ί <sub>ΟΗ</sub> = -80μΑ
Icc	V <sub>CC</sub> Current Drain		120	mA	
hι	Input Leakage		±10	μA	$0V \leq V_{IN} \leq V_{CC}$
<b>IOFL</b>	Output Leakage During Float		±10	μA	0.45V ≤ V <sub>OUT</sub> ≤ V <sub>CC</sub>

## **CAPACITANCE** ( $T_A = 25^{\circ}C; V_{CC} = GND = 0V$ )

Symbol	Parameter	Min.	Тур.	Max.	Unit	Test Conditions
CIN	Input Capacitance			10	рF	fc = 1MHz
C <sub>1/O</sub>	I/O Capacitance			20 ,	pF	Unmeasured pins returned to GND

## A.C. CHARACTERISTICS—PERIPHERAL (SLAVE) MODE

(8257:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 5\%$ , GND = 0V) (8257-5:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 10\%$ , GND = 0V)

#### 8080 Bus Parameters

## READ CYCLE

		8257		8257-5			
Symbol	Parameter	Min.	Max.	Min.	Max.	Unit	Test Conditions
T <sub>AR</sub>	Adr or CS↓ Setup to RD↓	0		0		ns	
T <sub>RA</sub>	Adr or $\overline{CS}^{\uparrow}$ Hold from $\overline{RD}^{\uparrow}$	0		0		ns	
T <sub>RD</sub>	Data Access from RD↓	0	300	0	220	ns	
T <sub>DF</sub>	DB→Float Delay from RD↑	20	150	20	120	ns	
T <sub>RR</sub>	RD Width	250		250		ns	

WRITE CYCLE

		82	8257					
Symbol	Parameter	Min.	Max.	Min.	Max.	Unit	Test Conditions	
TAW	Adr Setup to WR↓	20		20		ns		
TWA	Adr Hold from WR1	0		0		ns		
T <sub>DW</sub>	Data Setup to WR1	200		200		ns		
T <sub>WD</sub>	Data Hold from WR1	10		· 10		ns		
Tww	WR Width	200		200		ns		

**OTHER TIMING** 

		82	57	8257-5				
Symbol	Parameter	Min.	Max.	Min.	Max.	Unit	Test Conditions	
T <sub>RSTW</sub>	Reset Pulse Width	300		300		ns		
TRSTD	Power Supply† (V <sub>CC</sub> ) Setup to Reset	500		500		μs		
T,	Signal Rise Time		20		20	ns		
Tf	Signal Fall Time		20		20	ns		
T <sub>RSTS</sub>	Reset to First I/OWR	2		2		tcy		

## A.C. CHARACTERISTICS-DMA (MASTER) MODE

(8257:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 5^{\circ}$ , GND = 0V) (8257-5:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 10^{\circ}$ , GND = 0V)

### TIMING REQUIREMENTS

Symbol	Parameter	8257		8257-5		11-14
		Min.	Max.	Min.	Max.	Unit
Тсү	Cycle Time (Period)	0.320	4	0.320	4	μS
T,	Clock Active (High)	120	.8T <sub>CY</sub>	80	.8T <sub>CY</sub>	ns
T <sub>QS</sub>	DRQ1 Setup to CLK1 (SI, S4)	120		120		ns
Тон	DRQI Hold from HLDA1 <sup>[1]</sup>	0		0		ns
T <sub>HS</sub>	HLDA1 or ISetup to CLKI(SI, S4)	100		100 ·		ns
T <sub>RS</sub>	READY Setup Time to CLK1(S3, Sw)	30		30		ns
T <sub>RH</sub>	READY Hold Time from CLKI(S3, Sw)	30		30		ns

## A.C. CHARACTERISTICS-DMA (MASTER) MODE

(8257:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 5\%$ ,  $\dot{GND} = 0V$ ) (8257-5:  $T_A = 0^{\circ}C$  to 70°C,  $V_{CC} = 5.0V \pm 10\%$ , GND = 0V)

**TIMING RESPONSES** 

Symbol	Parameter	8257		8257-5		Unit
		Min.	Max.	Min.	Max.	Viiit
T <sub>DQ</sub>	HRQ† or ↓Delay from CLK† (SI, S4) (measured at 2.0V)		160		160	ns
T <sub>DQ1</sub>	HRQ† or ↓Delay from CLK† (SI, S4) (measured at 3.3V) <sup>[3]</sup>		250		250	ns
TAEL	AEN† Delay from CLK↓ (S1)		300		300	ns
TAET	AEN↓ Delay from CLK† (SI)		200		200	ns
TAEA	Adr (AB) (Active) Delay from AEN <sup>†</sup> (S1) <sup>[1]</sup>	20		20		ns
TFAAB	Adr (AB) (Active) Delay from CLK <sup>†</sup> (S1) <sup>[2]</sup>		250		250	ns
TAFAB	Adr (AB) (Float) Delay from CLK <sup>†</sup> (SI) <sup>[2]</sup>		150		150	ns
TASM	Adr (AB) (Stable) Delay from CLK <sup>†</sup> (S1) <sup>[2]</sup>		250		250	ns
T <sub>AH</sub>	Adr (AB) (Stable) Hold from CLK <sup>†</sup> (S1) <sup>[2]</sup>	T <sub>ASM</sub> -50		T <sub>ASM</sub> -50		ns
TAHR	Adr (AB) (Valid) Hold from $\overline{RD}$ (S1, SI) <sup>[1]</sup>	60		60		ns
TAHW	Adr (AB) (Valid) Hold from Wrt (S1, SI) <sup>[1]</sup>	300		300		ns
TFADB	Adr (DB) (Active) Delay from CLK <sup>†</sup> (S1) <sup>[2]</sup>		300		300	ns
	Adr (DB) (Float) Delay from CLK <sup>†</sup> (S2) <sup>[2]</sup>	T <sub>STT</sub> +20	250	T <sub>STT</sub> +20	170	ns
TASS	Adr (DB) Setup to Adr Stb1 (S1-S2)[1]	100		100		ns
TAHS	Adr (DB) (Valid) Hold from Adr Stb1 (S2)[1]	20		20		ns
TSTL	Adr Stb† Delay from CLK† (S1)		200		200	ns
TSTT	Adr Stb↓ Delay from CLK↑ (S2)		140		140	ns
T <sub>SW</sub>	Adr Stb Width (S1-S2) <sup>[1]</sup>	T <sub>CY</sub> - 100		T <sub>CY</sub> -100		ns
TASC	Rd↓ or Wr(Ext)↓ Delay from Adr Stb↓ (S2) <sup>[1]</sup>	70		70		ns
т <sub>DBC</sub>	RD↓ or WR(Ext)↓ Delay from Adr (DB) (Float) (S2) <sup>[1]</sup>	20		20		ns
T <sub>AK</sub>	DACK <sup>†</sup> or <sup>↓</sup> Delay from CLK <sup>↓</sup> (S2, S1) and TC/Mark <sup>†</sup> Delay from CLK <sup>†</sup> (S3) and TC/Mark <sup>↓</sup> Delay from CLK <sup>†</sup> (S4) <sup>[4]</sup>		250		250	ns
T <sub>DCL</sub>	$\overrightarrow{RD}$ or $\overrightarrow{Wr}(Ext)$ Delay from CLK <sup>†</sup> (S2) and Wr Delay from CLK <sup>†</sup> (S3) <sup>[2,5]</sup>		200		200	ns
T <sub>DCT</sub>	Rd <sup>†</sup> Delay from CLK↓ (S1, SI) and Wr <sup>†</sup> Delay from CLK <sup>†</sup> (S4) <sup>[2,6]</sup>		200		200	ns
TFAC	Rd or Wr (Active) from CLK <sup>↑</sup> (S1) <sup>[2]</sup>		300		300	ns
TAFC	Rd or Wr (Active) from CLK <sup>↑</sup> (S1) <sup>[2]</sup>		150		150	ns
TRWM	Rd Width (S2-S1 or SI) <sup>[1]</sup>	$2T_{CY}+T_{\theta}-50$		$2T_{CY}+T_{\theta}-50$		ns
T <sub>WWM</sub>	Wr Width (S3-S4) <sup>[1]</sup>	Т <sub>СҮ</sub> – 50		T <sub>CY</sub> -50		ns
TWWME	WR(Ext) Width (S2-S4) <sup>[1]</sup>	2T <sub>CY</sub> -50		2T <sub>CY</sub> -50		ns

NOTES:

1. Tracking Parameter.

2. Load = + 50 pF.

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